## CS 561, HW3

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Due: October 19th

- 1. Problem 4-6 (VLSI chip testing) This is a really good divide and conquer problem I left out of the last hw
- 2. Exercise 12.2-4 (Prof. Bunyan's property)
- 3. Exercise 12.4-2 ("Describe a binary search tree on n nodes such that the average depth...")
- 4. Problem 12-3 (Average Node Depth in Randomly Built Binary Search Tree)
- 5. Exercise 13.1-6 (Largest number of nodes with black-height k)
- 6. Exercise 13.3-1 ("In line 16 of RB-Insert ...")
- 7. Problem 13-3 (AVL Trees)
- 8. Problem 13-4 (Treaps)
- 9. HAFTs: A half-full tree (HAFT) is a rooted binary tree that is a useful data structure for designing self-healing networks. Let  $\ell$  be a positive integer. For  $\ell$  a power of 2, the complete tree with  $\ell$  leaf nodes is the unique haft with  $\ell$  leaf nodes. For  $\ell$  not a power of 2, a tree with  $\ell$  leaf nodes is a haft if and only if (1) the root node, r, has two children; (2) the left subtree of r is the root of a complete binary containing  $2^{\lfloor \log \ell \rfloor}$  leaf nodes; and (3) the right subtree of r is a haft. (Recall that a complete binary tree is one where every internal node has two children and every leaf node has the same depth)

Show the following by induction:

• For all positive  $\ell$ , there is a unique haft with  $\ell$  leaf nodes.

- Call the haft with  $\ell$  leaf nodes  $haft(\ell)$ . Then, the height of  $haft(\ell)$  is  $\lceil \log n \rceil$
- 10. Challenge: In the self-healing application of hafts, the leaf nodes are associated with actual machines in a network, and the internal nodes represent additional "router nodes" (a scarce resource). To merge a list of hafts,  $h_1, h_2, \ldots, h_x$  we want to create a single new haft, h, which contains as leaf nodes all the leaf nodes in  $h_1, h_2, \ldots, h_x$ , and adds the smallest number of new internal nodes as possible.
  - Show how you can merge a collection of x hafts, each of size no more than n, into a single big haft by adding no more than  $O(x \log n)$  internal nodes.

Hint: Think about how to set up a correspondence between binary numbers and hafts, and binary addition and haft merging.